Securing Flash and Solid State Drives

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Why Securing Flash

- Specific Requirements? Yes!!!
  - Confidentiality: an attacker can easily read data. Although wear-leveling can be considered a random permutation of the data block this is only obfuscation.
  - Integrity: an attacker can easily modify data blocks.
  - Access Control: NAND Flash does not provide any access control.

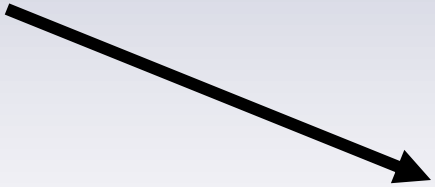
Why Security in Flash Storage?

3 Simple Reasons

- Storage for secrets with strong access control
  - Arbitrarily large memory space
  - Gated by access control
- Unobservable cryptographic processing of secrets
  - Processing unit “welded” to storage unit
  - “Closed”, controlled environment
- Custom logic for faster, more secure operations
  - Inexpensive implementation of modern cryptographic functions
  - Complex security operations are feasible

Securing Flash Storage

- Storage Security is not only about encryption:
  - Is a about:
    - Confidentiality
    - Integrity
    - Access Control
    - Key Management
    - Online and Offline
    - …..



Requires a platform capable of accommodating all these requirements.
Moreover in order to enable a storage security ecosystem a standardized platform is necessary.

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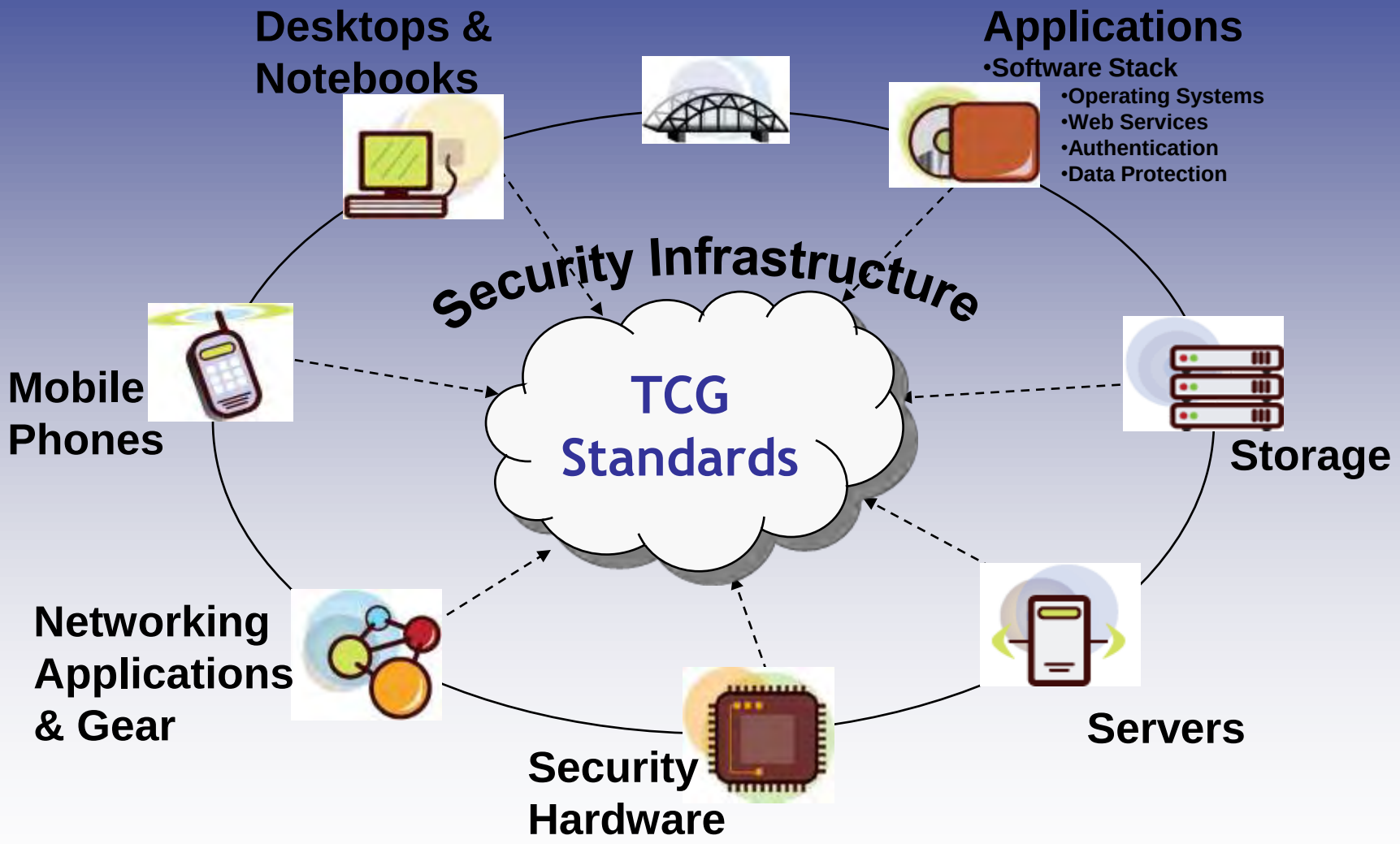
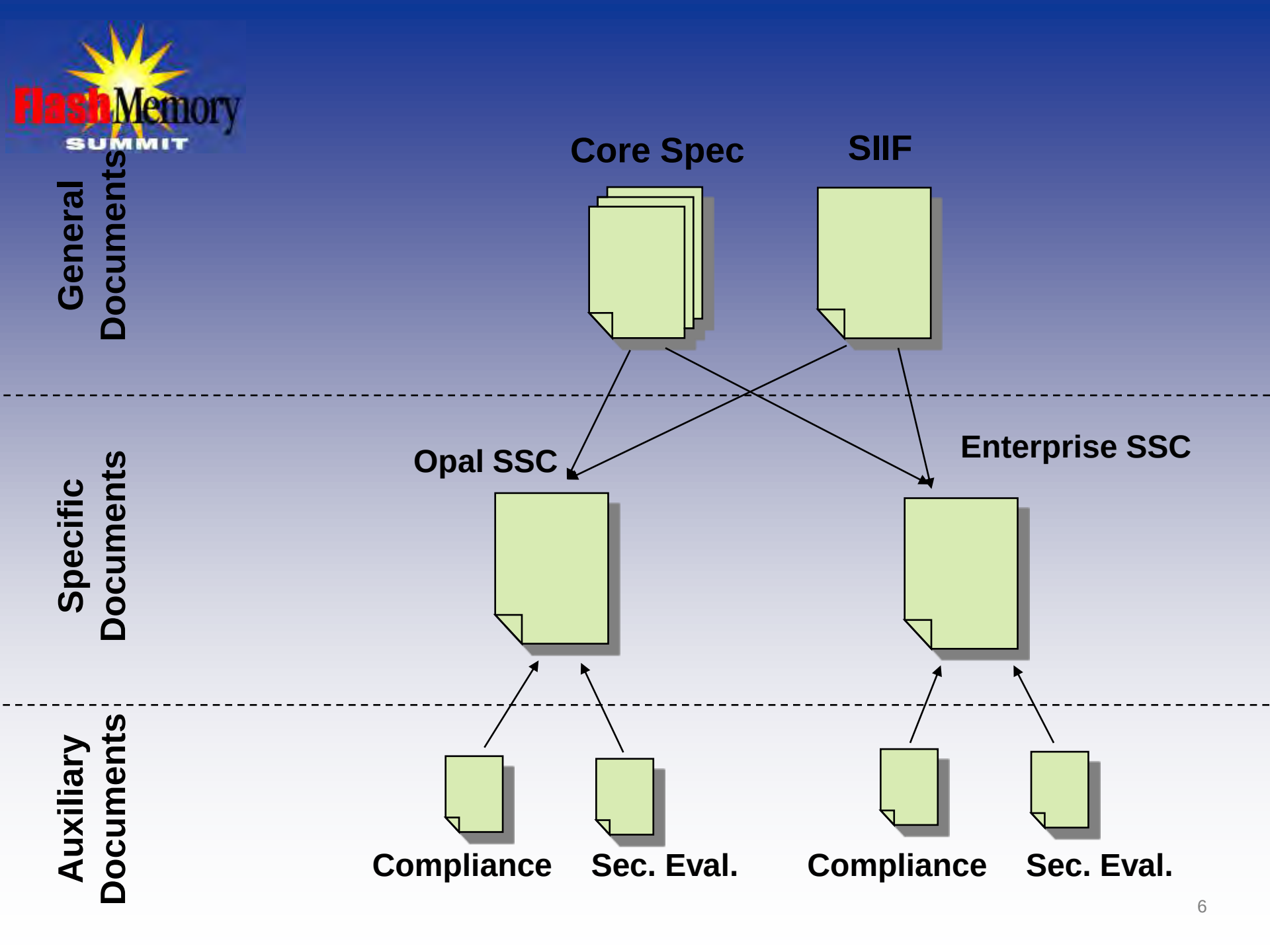
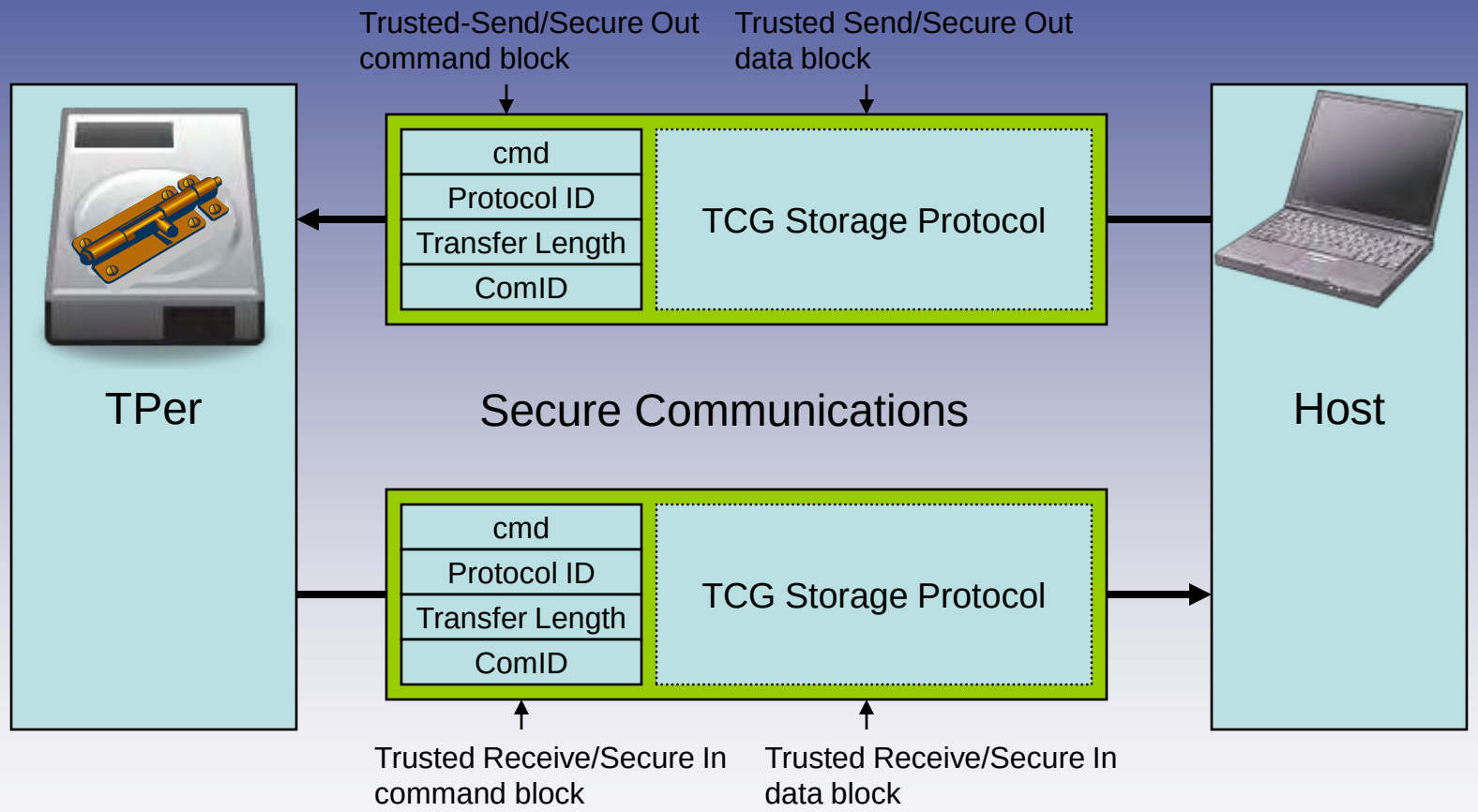


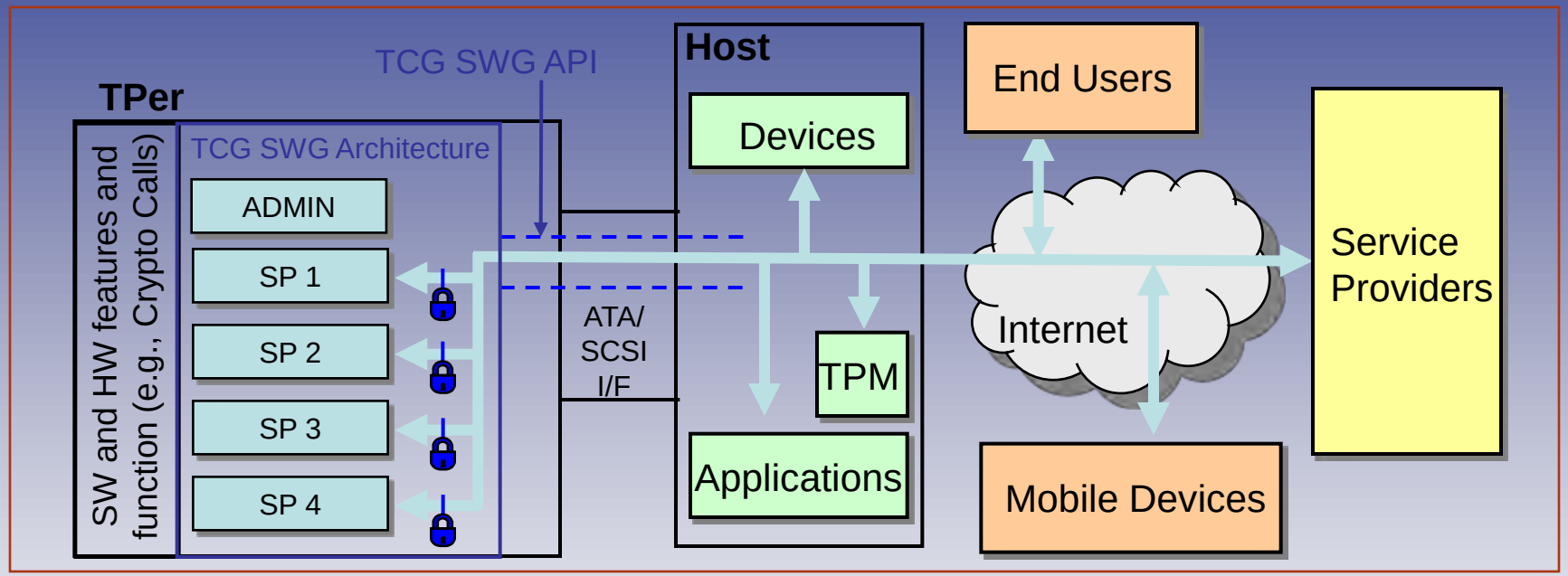

General Documents
Specific Documents
Auxiliary Documents
Core Spec
SIIF
Opal SSC
Enterprise SSC
Compliance
Sec. Eval.
Compliance
Sec. Eval.

Storage-to-Host Communications



ComID: allows TPer to identify caller of Trusted Receive/Secure In command

TCG Storage Work Group: Overview

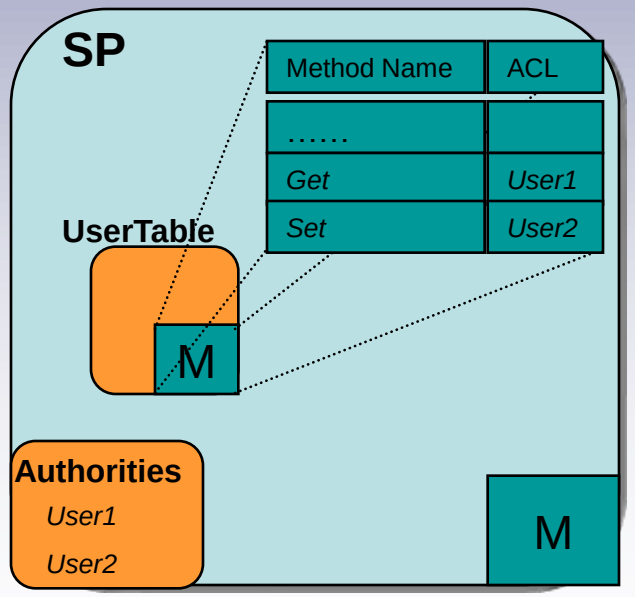


The host platform, applications, devices, local end users, and remote users/service providers can gain exclusive control of selected features of the storage device. This allows them to simultaneously and independently extend their trust boundary into the storage device or trusted peripheral (TPer)

TCG Storage Work Group: Architecture

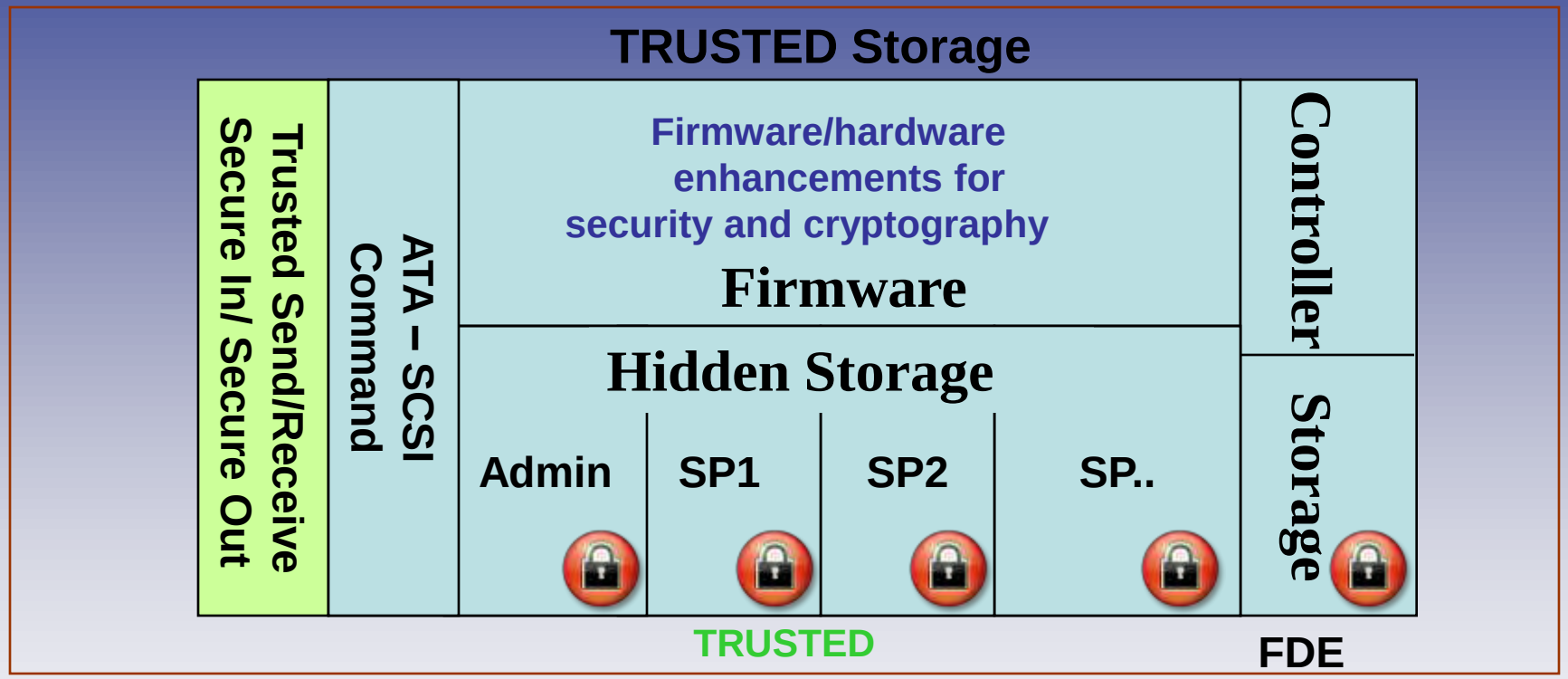
Storage Work Group specifications are intended to provide a comprehensive command architecture for putting selected features of storage devices under policy-driven access control.

- Features will be packaged into individual functionality containers called: "Security Providers" or SPs.



- Each SP is a "sand box" exclusively controlled by its owner. SP functionality is a combination of pre-defined functionality sets called SP Templates:
  - Base
  - Admin
  - Crypto
  - Log
  - Clock
  - Locking
- SPs are a collection of tables and methods that control the persistent trust state of the TPer.
  - Method invocation occurs under access control.
  - The SP has a list of authorities and their respective credentials for access control.

TCG Storage Work Group: Implementation Overview



- (Partitioned) Hidden Storage
- Security firmware/hardware
- Trusted Container Commands

Core Architecture Incarnations: Security Subsystem Classes

- The Core Specification defines a comprehensive set of security features, but not all are necessary to implement a security solution.
- An SSC defines such a subset addressed to a particular set of requirements/market.
- Currently under development:
  - NB Market / HDD Loss & Theft (Opal SSC)
  - Enterprise: Band Encryption and re-purposing (Enterprise SSC)

Enterprise SSC Overview

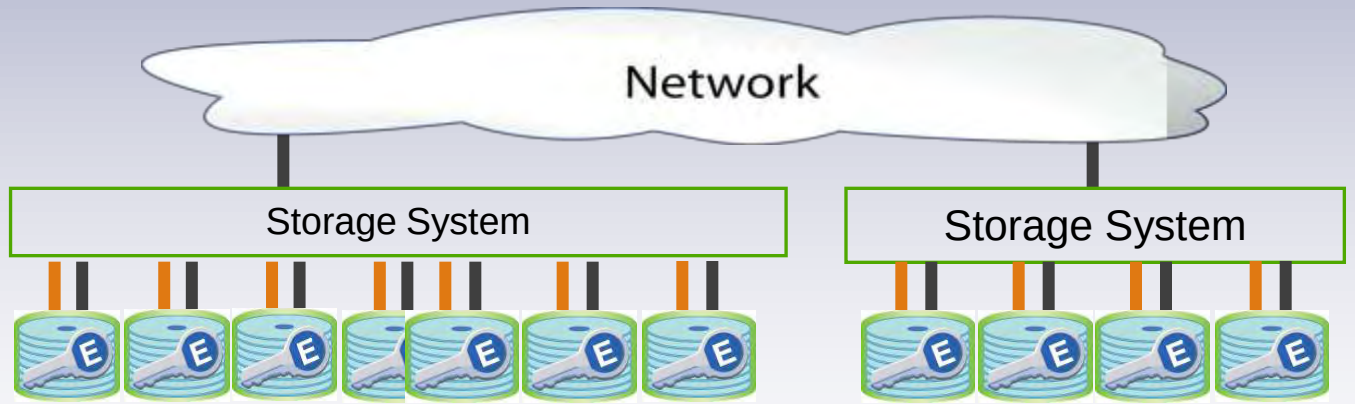
Encryption engine speed

Matches

Port's max speed

- Threat Model
  - Lost / Stolen Drives
- Features
  - Encryption
  - Drive Locking with PW access control
  - Encryption Ranges
  - (Fast Secure Erase)

Scales Linearly, Automatically



All data can be encrypted, with no performance degradation

Less need for data classification

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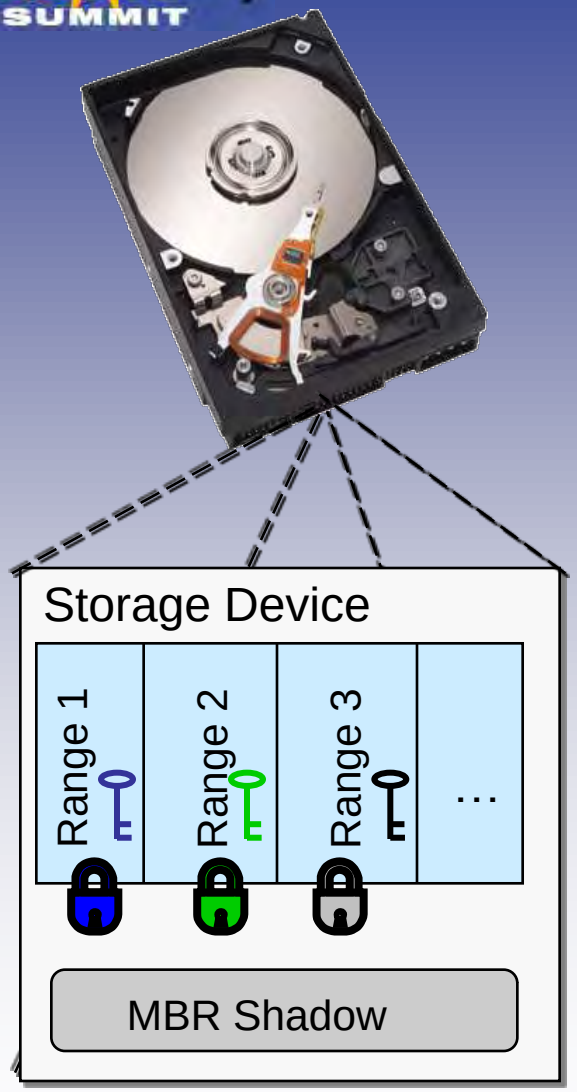
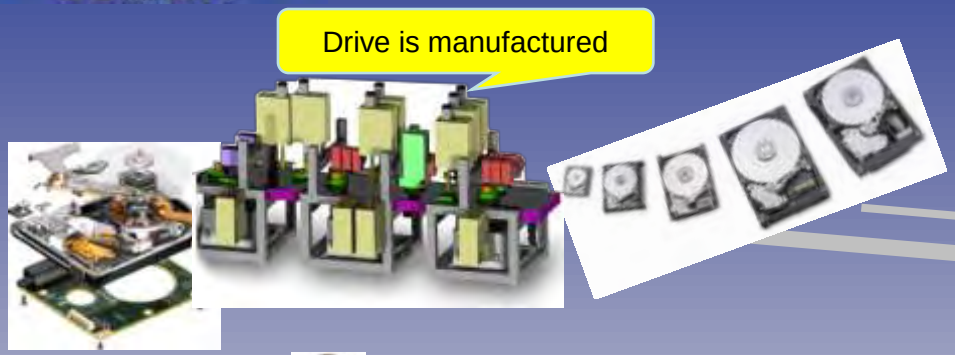
Opal SSC Overview

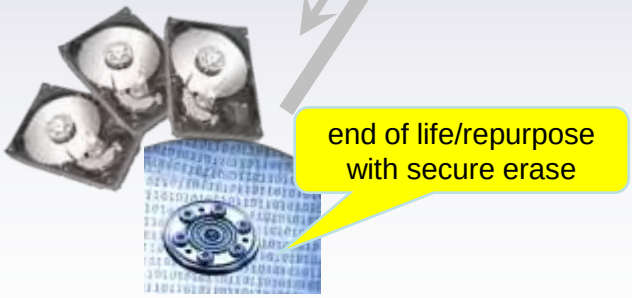


- Threat Model
  - Lost / Stolen Laptops
  - (Offline leakage of data)
- Features
  - Encryption
  - Drive Locking with PW access control
  - Encryption Ranges
  - MBR Shadowing (Pre-Boot)
  - (Fast Secure Erase)

- Very simple to use SSC addressing PC Client system needs.

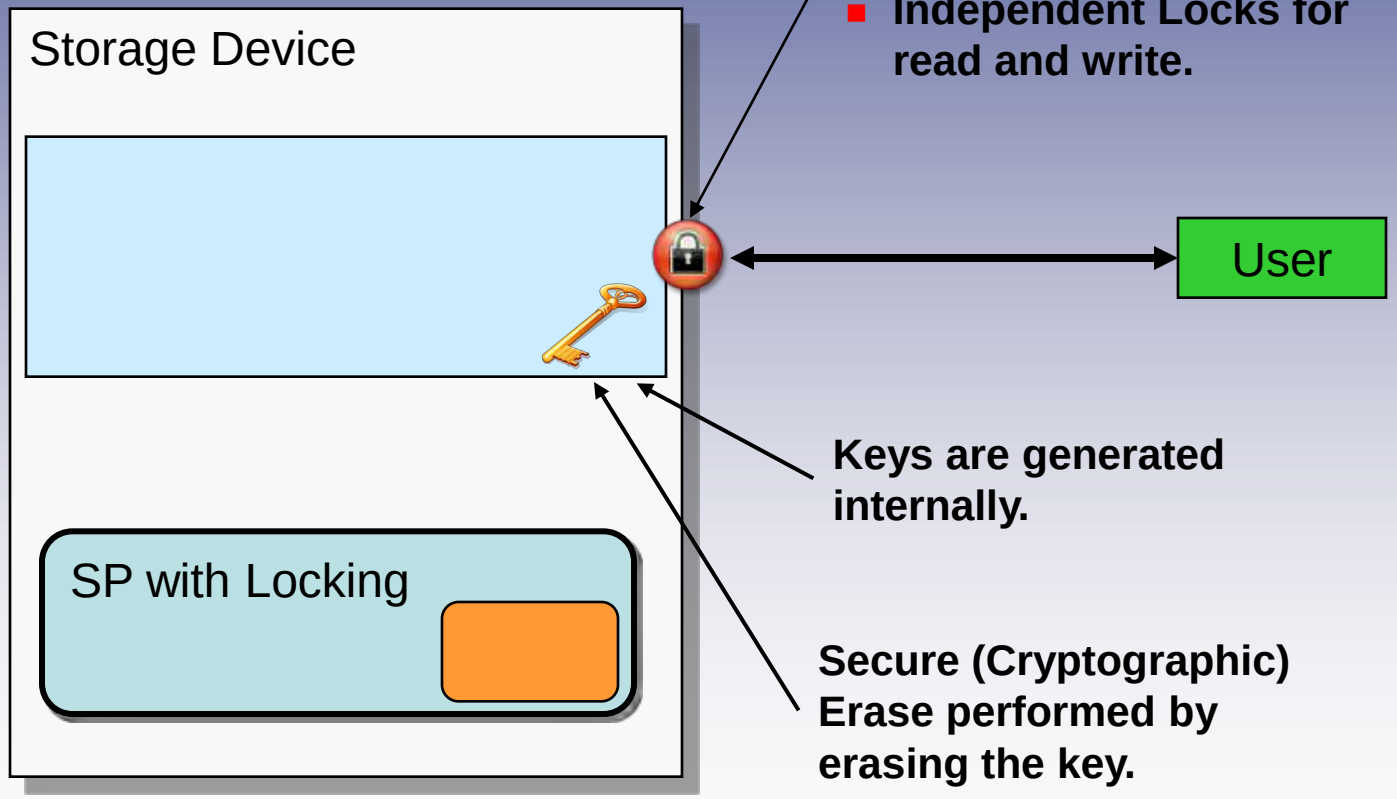
Example Life of an Opal HDD





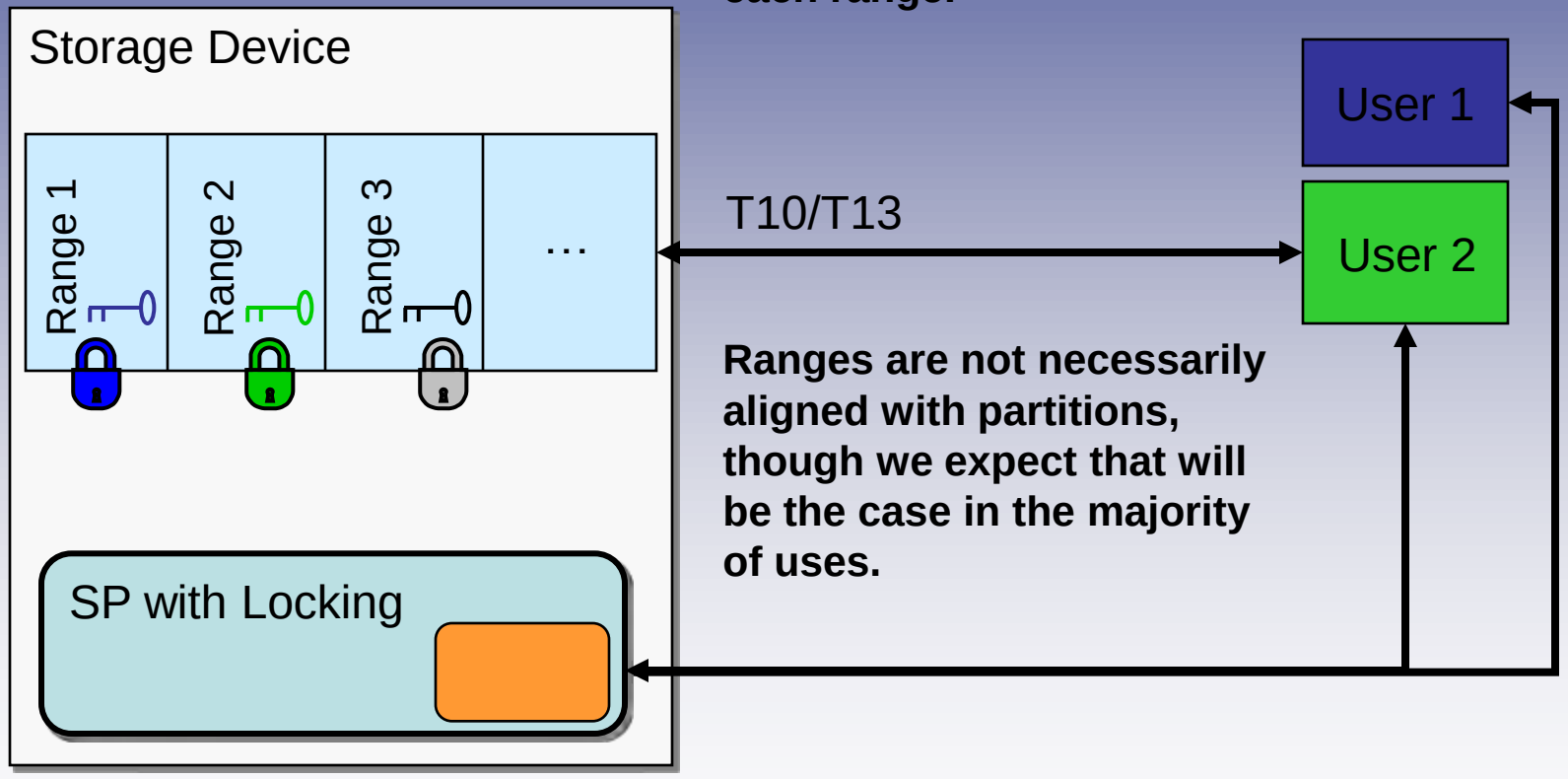
Encryption/Locking

- **Access control for locking and unlocking**
- **Independent Locks for read and write.**



Ranges

Independent encryption and access control for each range.

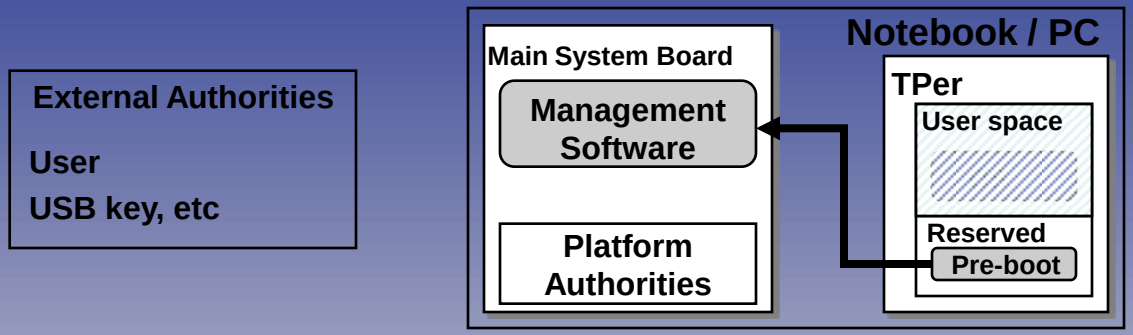


Ranges are not necessarily aligned with partitions, though we expect that will be the case in the majority of uses.

MBR Shadowing

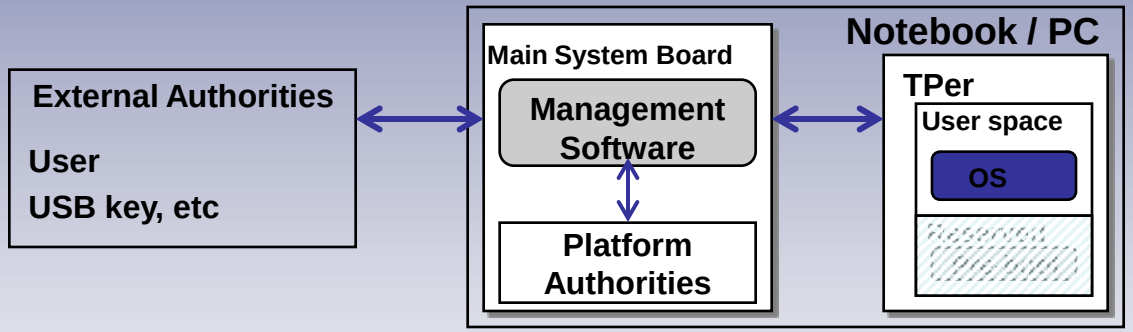
Initial Power-up

- When the system first requests the MBR, the HDD returns the pre-boot code (the MBR shadow).



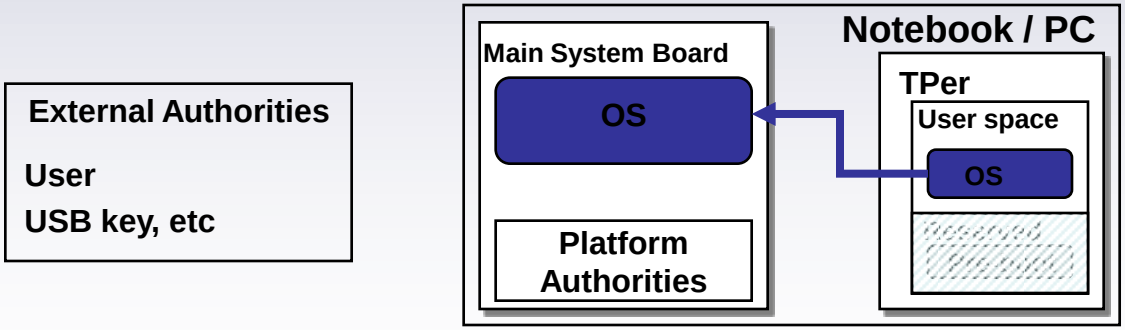
Authentication and Unlock

- The pre-boot code manages the authentication process with both internal and external authorities.
- After the appropriate authentications, the management software unlocks the regular user space.



Resume Normal Boot

- After the HDD is unlocked, the management software sends the system back to the boot process.
- The system's request for the MBR now returns the true MBR and the OS is loaded completing the boot process.



THANK YOU!

www.trustedcomputinggroup.org

Core Specification v2.0:

http://www.trustedcomputinggroup.org/resources/tcg_storage_architecture_core_specification_version_200_revision_100

Opal Specification v1.0:

http://www.trustedcomputinggroup.org/resources/tcg_storage_security_subsystem_class_opal_version_100_revision_200

Enterprise Specification v1.0:

http://www.trustedcomputinggroup.org/resources/storage_work_group_storage_security_subsystem-_class_enterpriseversion_10_revision_10