

Extending Memory Lifetime Using Coding

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- Problem:
	- o Cell erasures and memory wearout.
- **Proposed approach:** o Data self-repair to minimize number of erasures based on algebraic techniques.
- **Summary and outlook:**
	- o Overcome physical degradation by novel mathematical solutions.

Memory Lifetime and Write/Erase **Operations**

 When only one cell needs to be erased, the whole block needs to be reset.

- For SLC \sim 10⁶ writes,
- For MLC $\sim 10^4 10^5$ writes (serious)

 F lash Memory Sum H it Q $\mathsf{rTLC}~\simeq 10^3-10^4$ writes (more serious). Santa Clara, CA 33 and 2012 and 2013 and 2013 a

Memory Lifetime and Write/Erase **Operations**

\blacksquare Programming (write) error is very costly.

Wastes erase cycles.

Memory Lifetime and Write/Erase **Operations**

• Write is incremental step pulse programming a.k.a. "guess-and-verify".

\blacksquare Being cautious affects latency.

- **Idea 1: Allow for sloppy writes.**
	- o Improves latency.
	- o Not wasting erase/write budget.
	- o Reliability ?
- **If Idea 2: Figure out what was intended to be** written based on other cells.
	- o Overshot values stay intact (for the time being).
	- o Redundancy ?

Allow writing overshoot, a.k.a. sloppy writes.

Key: unidirectional error correction scheme.

Varshamov-Tenengolts codes [1]:

$$
\sum_{i=1}^{n} ix_i \equiv a \bmod (n+1)
$$

- o *n* is block size,
- o *xi* is value in cell *i*,
- o *a* is arbitrary integer.
- VT code corrects one unidirectional error.

[1] R. R. Vashamov and G. M. Tenengolts, 1965.

Proposed scheme:

$$
\sum_{i=1}^{n} i x_i \equiv a_1 \mod p
$$

$$
\sum_{i=1}^{n} i^2 x_i \equiv a_2 \mod p
$$

$$
\sum_{i=1}^{n} i^{2k} x_i \equiv a_{2k} \bmod p
$$

...

Parameters:

- o *k* is target error correction
- o*n* is block size
- o *xi* is value in cell *i*,
- *a₁…a_{2k} are* arbitrary integers.

o *p* is some prime, *p* > *n*

Ξ Guaranteed to correct *k* unidirectional errors.

- **Encoding:**
- 1.Compute congruency contribution from data.
- 2.Add values in anchors for overall congruency.
	- o With careful indexing, redundancy is minimal.
	- oSystematic construction.

Example: $\textbf{target:} \, \, \sum \, \textsf{i} \, \, \textsf{x}_{\textsf{i}} \,$ = 0 mod p

index 1 2 3 4 5 6 7 8 9 10 values **1 1 1 1 0 1 1 0 0 1**

- **Decoding**
	- 1. Test if congruency constraints are violated.
	- 2. Solve equations to figure out the correct values. \overline{O} Computations can be efficiently implemented.
- **Example: target:** \sum i x_i = 0 mod 11

index 1 2 3 4 5 6 7 8 9 10 values **1 1 1 1 0 1 1 0 1 1**

- **Decoding**
	- 1. Test if congruency constraints are violated.
	- 2. Solve equations to figure out the correct values.
		- oComputations can be efficiently implemented.
- **Example:** $\textbf{computed:} \; \sum i \; \text{x}_\text{i} = 9 \; \text{mod} \; 11$

index 1 2 3 4 5 6 7 8 9 10 values **1 1 1 1 0 1 1 0 1 1**

- **Data self-repair can improve write latency and** extend memory lifetime.
- **Efficient methods are developed based on** number-theoretic ideas.
	- o Very low redundancy
	- o Implementable algorithms o For SLC/MLC/TLC
- Rich opportunity to develop new data correction algorithms and methodologies tailored for Flash.

- Thank you for your attention!
- For more information

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