

# Making ECC Work For Flash Part II

### **Optimizing Flash ECC and RAID**

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Flash Memory Summit 2014

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- A discussion of UBER (definitions)
- The JEDEC specs and what they mean
- UBER/NRRE and RAID
- Failure targets
- DNR ECC
- New RAID codes



### **UBER: Uncorrectable Bit Errors**

#### The "other" component of reliability

- UBER is when there are more bit errors than the ECC can correct
  - For example, if the sector ECC can correct 50 bits, but there are more than 50 bits in error
- One component of non-recoverable read errors (NRRE)
  - 2 outcomes of an NRRE event:
    - The ECC detects the error count is too large, and declares the sector lost
    - The ECC blissfully applies the correction and produces an incorrect value (miscorrection)
      - We tend to add CRC to detect these events and turn them into NRRE events
- I'll use NRRE going forward in this analysis

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### **Non-Recoverable Read Errors**

- NRRE events contribute to data loss
  - Impact depends on the system architecture
  - Loss is at least a sector worth of bits
- NRRE is specified as an interval: e.g. < 1 in  $10^{14}$  bits
- Or as a rate: e.g.  $\leq 10^{-14}$  per bit
- 10<sup>14</sup> bits seems really large
  - But there are  $0.08 \times 10^{14}$  bits in a terabyte!



# **Alternate NRRE Specifications 1**

Some alternate approaches

- Express as rate per TB transferred
  - Nice for computing from data moved
  - NRRE/TB = error\_interval/8 x  $10^{12}$
- Express as sector failure probability per operation (sector read)
  - More accurate, since we lose a sector on an NRRE event, not a bit
  - psfail = sector\_bits/error\_interval



Some typical specifications (assume 1kB sectors)

	Consumer HDD	Enterprise SSD
Typical NRRE Spec (b)	1e14	1e17
NRRE/TB	8%	8e-5
psfail	8.2e-11	8.2e-14

#### Observation: the new metrics are more informative

Which may explain why the vendors haven't adopted them

emorv

FIa



### **NRRE Specs and Data Loss**

- SSDs need much tighter NRRE specs than HDDs
- SSD industry has set specs based on HDDs
  - Unfortunately, industry hasn't quite noticed the need for improvement
- We can estimate data loss rates from specs and workload
  - Workload will be small block random IO (why we use SSDs)



### **NRRE Specs and Data Loss**

#### Simple to estimate

- psfail = sector\_bits/error\_interval
- Sector\_Ops/Y = 3,600\*8,760\*IOPS\*sectors\_per\_IO
- Mean Y/Sector Loss = 1/(Sector\_Ops/Y \* psfail)
- Can add duty cycle effects, but these are small – R/W typically 70/30
  - Active duty cycle ~80% enterprise, ~20% consumer



#### SSDs running at spec are at high risk of data loss

	Con HDD	Con SSD	Ent HDD	Ent SSD
IOPS (4kB)	100	10,000	350	150,000
Sector Ops/Y	1.3e10	1.3e12	4.4e10	1.9e13
NRRE Interval (bits)	1e15	1e16	1e16	1e17
psfail	8.8e-12	8.8e-13	8.2e-14	1.6e-15
Mean Y/Sector Loss	10	1	28	0.6
MTTDL (Hours)	85k	8.5k	242k	5.5k
Scaled NRRE Interval		1e17		5e18

- Both consumer and enterprise SSD NRRE specs are too loose
  - Duty cycle effects impacts consumer more than enterprise
  - Will see the effects more at high PE cycle counts
- We need tighter specs!
  - JEDEC specs (JESD218) are 1e15 and 1e16

Shorter than the MTBF!



#### Protecting against device loss and sector loss

- When we use the term RAID we refer to an erasure correcting code that protects against unit loss
  - Can be hardware or software based
- NRRE impacts reliability during rebuilds
  - If there is no parity left, a sector loss becomes a data loss event
  - Occurs when rebuilding a first failure in RAID 5
  - Occurs when rebuilding a second failure in RAID 6
  - Usually higher NRRE probability than a further unit failure during rebuild
    - Rebuild windows are short
    - Declustering parity doesn't help sector loss



### DNR ECC ("Do Not Resuscitate")

# Allow the NRRE (sector loss rate) at the SSD to be much greater, and let the RAID layer reconstruct the data





#### We let sectors fail at a higher rate with DNR ECC

- Failure (at the flash layer) is acceptable in RAID with larger limits than solo devices permit
- System can be optimized by adjusting the correction at each level
- No need to try so hard at the flash layer
  - DNR we deliberately set a higher failure rate target at the component level
  - Improves flash efficiency, simplifies encode/decode
    - Need to correct fewer errors
    - Makes the components more testable



#### How to create data loss targets for a system

- Failure events should be expressed per unit time
  - This is how the customer experiences events
    - Not per byte, or per IO
- Program based targets
  - Look at the behavior of an entire field population
    - Helps for modeling warranty costs
    - Also helps with program financial targets
- Inputs
  - Install base
    - Unit ships per year, field lifetime, program lifetime
  - Usage characteristics
    - Total data operations, total data transferred
  - Failure tolerance
    - Depends on the failure type
    - Is it a warranty event, loss of availability, loss of data or customer neardeath experience?



#### Precision is highly over rated here

- We need only compute first order terms!
- Why?
- Our assumptions are errors are independent of each other and of time
  - These are rarely true
    - (Well, essentially not at all with NAND...)
- The biggest deviations will be these assumptions
- So first order is good enough
  - Still a good idea to verify which terms are second order
- Thus, we can compute from binomials
  - Easy to do in a spreadsheet too!



### System Data Loss Targets

Program Design	ı	V	alue	N	otes		
Field lifetime Y	Field lifetime Y 5		5		Typical		
Mean field unit	Mean field units		1,000,000		Assume a successful program		
Units/Array		1	10		RAID span		
IO size (kB)		4		A	Assume transaction processing		
Total field IOs		3.	.15e19	A	Assume 50,000 IOPS/unit		
Arrays/field		100,000					
Program Loss Targets V			Value	Notes			
Data Loss Event	s/program	۱	1	For the entire program			
Target Prob data loss/array/Y2e-6			2e-6	A	ssume a successfu	ıl program	
SSD Device	Value	SSD Device		Value	SSD Device	Value	
IOPS	50,000	AFR		0.5%	ECC ovh bits	924	
Capacity (TB)	1	NRRE		1e16	Total bits/sect	9,212	
Sector kB	1	ECC Corr bits		66	Data Eff.	89%	
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### **Cumulative Binomial**

#### Useful for estimating failures

- cumbinomial(fails,trials,errorrate)
  - Fails is the number of failures
  - Trials is the total number of events (ops, bits, etc.)
  - errorrate is the failure rate per trial (e.g. ber)
- This is the cumulative binomial distribution
   In Excel, use the Binom.Dist function as:
  - 1-Binom.Dist(fails,trials,errorrate,TRUE)
  - Beware sometimes this runs out of precision when it shouldn't



#### Probability a RAID array has lost 1 unit in a year

- P1fail is the probability there is one failure in an array
- Probability an array is down 1 unit:
  - P1fail/Y = 1-binomial(0,arraysize,AFR) = 4.9% here
  - Not surprising:
    - 0.5% AFR \* 10 units = 5%
    - 0.5% AFR = 1.75MH MTBF



The required psfail which meets the system target

- psfail1f is the sector failure rate with 1 unit failure - We have a 1TB SSD and1kB sectors here
- psfail1f needed to meet array data loss target with 1 unit failure (RAID 5)
  - psfail1f = TgtDataLoss/Y / (sectorsread\*P1fail/Y)
  - Sectorsread = (1TB/1kB)\*(10 1)
  - psfail1f = 4.55e-15
  - (NRRE1f = 4.87E-19 is the equivalent NRRE to psfail1f)



## **Device Sector Fail and NRRE**

#### Our device is out of spec for the system

- Recall our consumer grade SSD had NRRE 1e16
- Which has psfail = 8.8e-13 but we need 4.55e-15!
- So, this device doesn't work here as specified
  - (No surprise, it's a consumer device)
  - But the enterprise drive at 1e-17 won't work either
- To continue, we will increase the ECC bits
  - Alternative is to limit the ber
  - Shouldn't change the answer much (either way it's a change to the SSD)



# **Computing NRRE Targets**

#### Get the raw bit error rate from the NRRE

- We can compute the raw ber from the psfail spec and ECC
  - 1. Assume BCH 66 code on 1kB
    - Corrects 66 bit errors out of 1,024 data bytes
    - Requires 924 check bits
    - sectorbits = databits + checkbits + metadata ~ 9,212
  - 2. psfail = (1 cumbinomial(66,sectorbits,ber))/sectorbits
  - 3. Invert by iteration to solve for ber
  - 4. Here: ber = 2.65e-3
  - 5. To meet system target need @ ber 2.65e-3 need 75 bits
    - 9,338 sectorbits
    - Hint: you can use Goal Seek in Excel to quickly iterate to find the ber

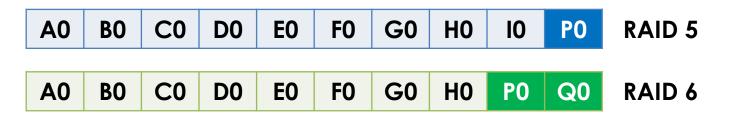
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### RAID 5 vs. RAID 6

#### Parity Groups

- Parity group (pgroup)
  - A collection of sectors that form an independent ECC set
  - In RAID 5 and 6 it's one sector from each unit
- RAID 6 has 2 parities per group
  - Can correct 1 sector/group after 1 unit failure
- Parity groups/array = 1e9 here
  - =TB/unit \* units/array \*1e9 / (sectors/group \* sector\_kB)
- Our arrays here chosen as 10 sectors/pgroup





Parity group failure on rebuild

- Sector data efficiency 0.9 RAID5, 0.8 RAID 6
  - (data sectors per group)/(total sectors per group)
- sparity is sector parities per group available after 1 unit failure
  - 0 for RAID 5, 1 for RAID 6
- Prob that a group fails to rebuild: pgroupfail = 1-cumbinomial(1+sparity,sectors/pgroup-1,psfail)



### RAID 5 vs. RAID 6

#### Rebuild failure and results

- Prob rebuild failure/array (multiple parity groups/array) prebuildfail = pgroupfail \* pgroup/array
- Absolute probability of array failure/year parrayfail = prebuildfail \* P1fail/Y
- Then, adjust the ECC correction bits to compute psfail until the parrayfail <= data\_loss-target
  - Since this is an integer, Goal Seek in Excel doesn't do as well



### DNR Results for RAID 5 vs RAID 6

RAID Type	RAID 5	RAID 6
sparity/pgroup	0	1
sectors/pgroup	10	10
pgroup/array	1e9	1e9
RAID data efficiency	0.90	0.80

Failure computations		
parrayfail	2.00e-6	2.00e-6
psfail	4.55e-15	3.37e-8
ECC corr bits	75	55
Sector efficiency	0.88	0.90

#### Not the answer the judges were looking for!

Net data efficiency

0.79 0.72 Steven Hetzler, IBM

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#### Sometimes you overpay for RAID 6 protection

- RAID 6 DNR doesn't increase the efficiency
- RAID 6 has 1 sector parity per parity group
  - These double as second unit failure protection
- What we need is a more efficient class of RAID
- What about parities designed for sector loss?
- fpof first point of failure
  - The minimum number of losses that cause a RAID failure



#### Optimized for both device and sector protection

- New RAID codes designed for this very problem
  - (I know, I was there at the time)
  - Parity group is now multiple sectors from each device (columns)

<b>A</b> 0	BO	<b>C0</b>	D0	EO	<b>P0</b>
<b>A</b> 1	<b>B</b> 1	C1	D1	E1	<b>P1</b>
A2	B2	C2	D2	E2	<b>P2</b>
A3	<b>B</b> 3	C3	D3	E3	<b>P3</b>
A4	<b>B4</b>	C4	q <sub>a</sub>	q <sub>b</sub>	P4

P0 is row 0 parity (Example with 6 units)
P1 is row 1 parity
P2 is row 2 parity
P3 is row 3 parity
P4 is row 4 parity, q<sub>a</sub>, q<sub>b</sub> group parities

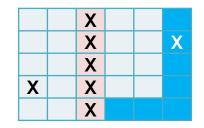
- Unit loss protection via row parities P
- Floating sector loss protection via group parities q
  - The q can be placed anywhere in the parity group
  - They are invoked only after more than 1 sector in a row is lost



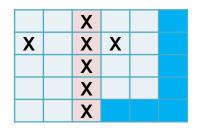
### RAID 5++

### • RAID 5++ in most cases stronger than RAID 6

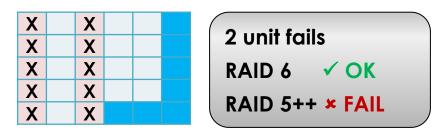
- Consider rebuild (1 unit fail)
- -RAID 6:
  - Correct all 1 sector fail/row
  - Correct 0 2 sector fail/row
  - fpof 2 sectors + 1 unit
- RAID 5++ :
  - Correct 2 1 sector fail/row
  - Correct 1 2 sector fail/row
  - fpof 3 sectors + 1 unit



2 rows with 2 fails RAID 6 ✓ OK RAID 5++ ✓ OK



1 row with 3 fails RAID 6 **\* FAIL** RAID 5++ **√** OK



- RAID 5++ is stronger to sector failure on rebuild
- RAID5++ is weaker to unit fails
  - Mitigated by short rebuild time



# DNR Results for RAID 5+ and RAID 5++

RAID Type	RAID 5	RAID 6	RAID 5+	RAID 5++
sparity/pgroup	0	1	1	2
sectors/pgroup	10	10	160	1,280
pgroup/Array	1e9	1e9	6.25e7	7.81e6
RAID data efficiency	0.90	0.80	0.89	0.90

Failure computations						
parrayfail	:	2.00e-6	2.00e-6	2.00e-6	2.00e-6	
psfail	4.	.55e-15	3.37e-8	7.2e-9	2.47e-7	
ECC corr bits		75	55	56	52	
Sector efficiency			0.90	0.90	0.91	
	We have a winner					
Net data efficiency		0.79	0.72	0.81	0.82	

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### **DNR Results with PMDS Codes**

#### RAID 5++ makes DNR ECC cost effective

- Efficiency is increased by letting the NRRE (psfail) increase – Up to 3% more efficient in this example
- May not sound like much, but worthwhile
  - Goes straight to margin
    - What else would you do for 3 margin points?
  - Can also be used to increase yields
  - May save cost in ECC decoders
  - Allows use of consumer parts in enterprise applications
- This was just a simple example, we may be able to do better with other configurations
- If you need dual failure protection, there are PMDS codes for those as well
  - If 2<sup>nd</sup> parity is protecting against a second unit failure, it's not available for sector loss protection
    - I have shown you how to do the math



### Summary

- Showed that UBER/NRRE specs for SSDs are inadequate – Time to data loss spec should be similar/better than HDD
- Have shown how to compute system reliability targets
- DNR ECC can achieve higher data efficiency
  - Allowing higher sector failure rates (NRRE) improves system cost
- PMDS codes such as RAID 5++ make DNR economical
- I will post a spreadsheet for downloading on my blog at smorgastor.DrHetzler.com

- Shows the details of the calculations for the interested student

• If you want to see actual data on SDD bit error rates and causes, attend my Tutorial T1 at 8:30 on Wed 8/6