

Building a Flash Optimized Storage System

Software Approach

Saumyabrata Bandyopadhyay and Ravi Dronamraju

Flash Memory Summit 2015 Santa Clara, CA

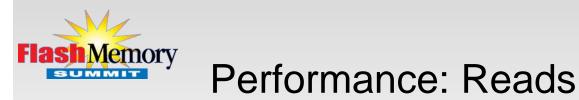


Flash : The change agent in Storage Software Design

- Performance
 - Enables 4-5 times improvement in throughput at low latency
 - Sustained read and write rate
 - Scale up vs. Scale out
- Physical Limits and Goals of Media
 - Finite life span of the media
 - Every potential reuse of the block result in media life loss
 - Better power usage
 - Smaller Form Factor and Weight
- Cost
 - Higher upfront cost
 - Should integrate seamlessly with lower endurance Flash
 - Offset with better TCO

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Techniques:

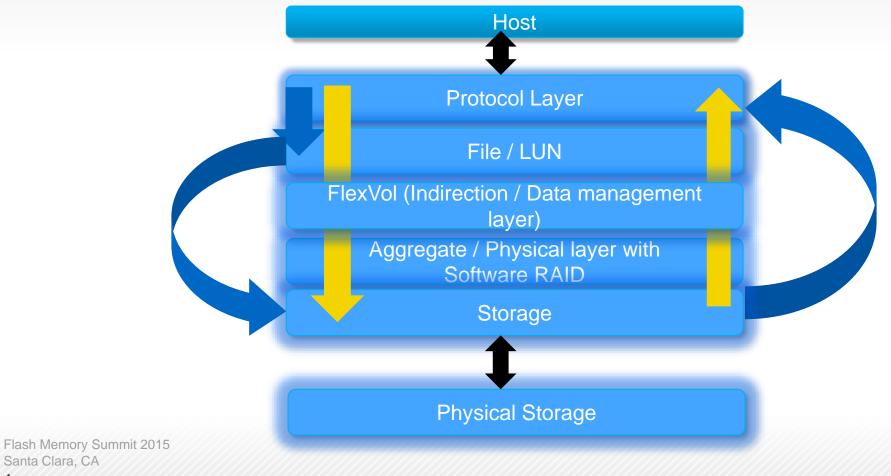
- IO path with few delay centers
- Use all the channels on the SSD
- Different priority for
 - Data being read by host vs. system
 - Data being read vs. written
- Software needs to be able to take advantage of all the CPU
 - MP safe data structures

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Optimized read path to take advantage of device latencies.



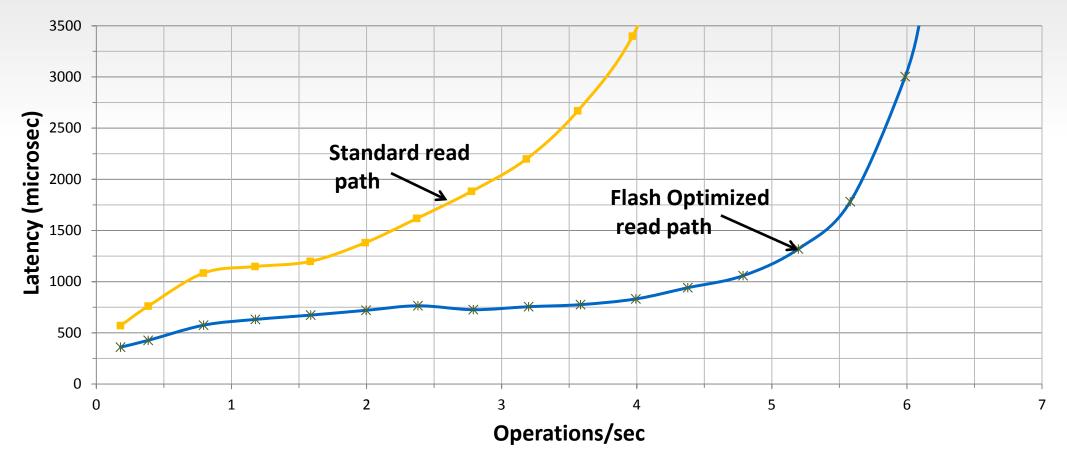


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Data ONTAP Read Path for Flash: Impact on performance



4x Improvement on Random Read Operations for the same latency





Performance: Writes

Techniques:

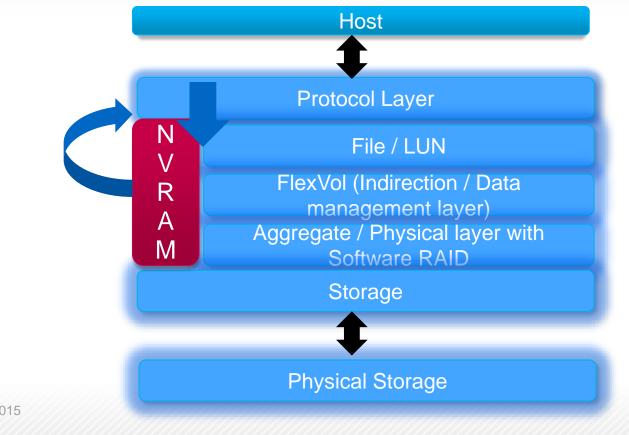
- Journal changes or log the operation to persistent lowest latency media
 - NVRAM/NVMem or battery backed DRAM
- Perform CPU intensive operations and metadata updates in the background instead of the IO path
- Writes to flash batched and sequential to reduce WAF
- Frees / overwrites batched to reduce random metadata updates





Data ONTAP IO Path: Writes

- Writes logged to NVRAM
- Data pushed to persistent storage periodically in the background



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Performance Trade Off: Scale out vs. Scale up

Performance / Application requirements dictate architecture trade off

Scale Up

- No fixed Storage to DRAM and CPU ratio
- Metadata must be searchable on disk for performance
- Applications with Small working sets but large storage footprint

Scale Out

- Fixed Storage to DRAM and CPU ratio
- Metadata normally in memory
- Performance scales with nodes
- Performance sensitive applications

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Data ONTAP: Scale Up and Scale out

- We need both
- Cost, performance and data center footprint dictate scale up vs. scale out
- Ideal performance when DRAM to Storage ratio around 1%
 - Not every workload / Application requires the same level of performance
- Scale up
 - Metadata for the working set fits in DRAM
 - Applications can tolerate occasional latency spikes during workload shifts
- Scale out
 - Consistent latency
 - Performance scales with storage
 - Large working sets

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Flash Physical limits: Layout

- Log structured layout
 - No writes in place
 - Results in fragmentation addressed using other means
- Indirection layer to address physical blocks
 - Ability to defrag & create clean segments
 - Reduces random overwrite, which in turn reduces WA
 - Allows for additional functionality that reduces the overall TCO
 - deduplication
 - cloning
 - thin provisioning
- Log On Log
 - Mitigated by writing sequentially to large segments using a single stream
 - Data segregation with multi-stream capabilities on next version SSD firmware





Data ONTAP – Layout Engine

- Write Anywhere File Layout (WAFL)
 - Log structured file system
 - Already optimized for write
 - NVRAM for staging writes
- Metadata searchable both in-memory as well as persistent storage
- Metadata updates optimized by batching updates
- FlexVols provide an indirection layer
- Generic File system instead of optimized for SAN
- Format that supports Storage Efficiency features
- Media aware physical layout engine allows for support of different types of persistent media
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Cost

Reduce the TCO by enabling Storage efficiency features

- Log structured file layout
 - Create space efficient snapshots with very little overhead
- Indirection layer allows for Storage efficiency features like
 - Deduplication
 - Clones
 - Thin provisioning
 - Writeable snapshots
 - Segment cleaning / Defrag which reduces the Write Amp on the SSDs
 - Allowing for reduction of Overprovisioning on the SSDs

Format that supports compression at various levels

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Conclusion

- Software design requires efficient use of CPU and fewer delay centers
- Software needs to customized for media writes
- Layout and format is important to reduce TCO by enabling storage efficiency features
- Application performance requirements dictate architecture

