



 Increasing Flash Throughput for Big Data Applications (Data Management Track)

Flash Memory Summit 2015

## Industry Context

- Addressing the challenge
- A proposed solution
- Review of the Benefits



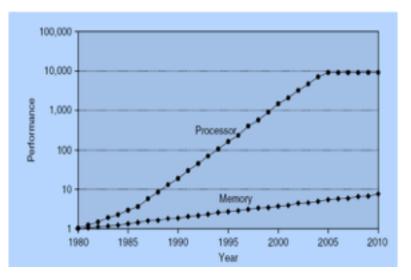
- Data sizes for data under management are monotonically increasing
  - Who wants less data?
- Our appetite for analysis is monotonically increasing
  - Do you think, or do you know?
  - Trend toward evidence-based management
- Our appetite for speed is monotonically increasing
  - Who wants questions answered more slowly?
  - Hence the industry interest in in-memory data management systems
- Our overall ability to manage complexity is <u>not</u> increasing



## **Processor / Memory Context**

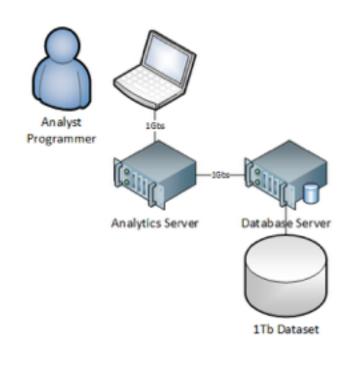
- Processor speeds are limited
- Processor core density has been increasing at a healthy rate
- Memory density is increasing (but at a lower rate than core density)!
- Therefore, the memory/core ratio is going in the wrong direction!
- We haven't significantly changed the memory/storage hierarchies for decades
  - Interconnects are getting faster as fast as memory access?
  - memory access is slow
  - caches are fast!

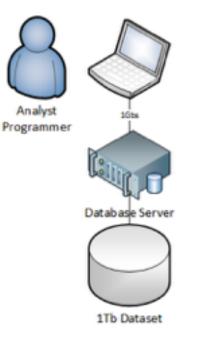
Flash Memory



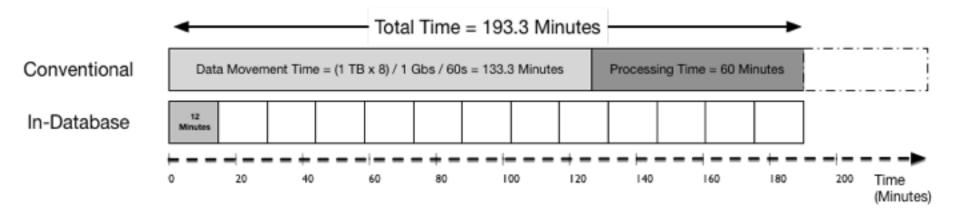
Hennessy, John L.; Patterson, David A. (2011-10-07). Computer Architecture: A Quantitative Approach (The











#### Moving data is the most time-consuming activity, reducing data movement is key. I/O is the enemy

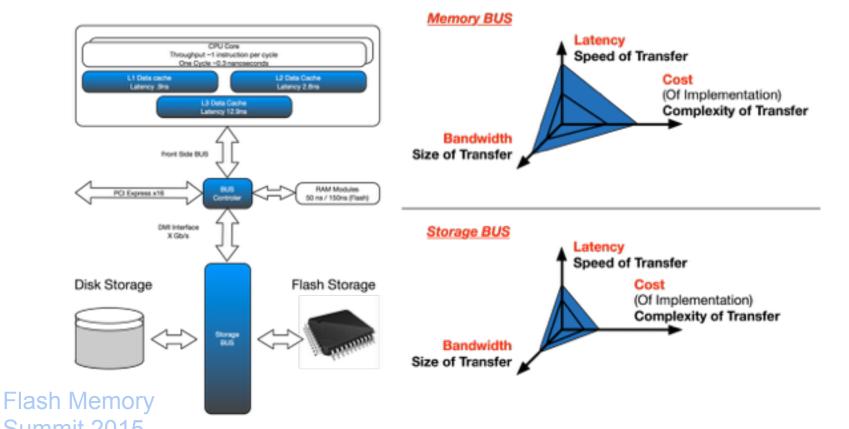


## The Memory Hierarchy in Human Terms

(.3ns = 1s)

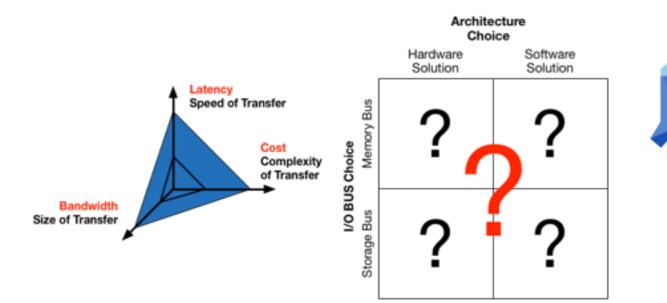
Event	Latency	Scaled
1 CPU Cycle	0.3 ns	1 s
Level 1 Cache Access	0.9 ns	3 s
Level 2 Cache Access	2.8 ns	9 s
Level 3 Cache Access	12.9 ns	43 s
Main Memory Access (DRAM, from CPU)	50.0 ns	3 min
Memory over Ethernet	3.2 µs	3.2 hours
CPU Context State Transfer	6.0 µs	6.0 hours
Flash SSD (PCI-e)	4.7 ms	5 months
Rotational disk I/O	1-10 ms	1-12 months
Internet: San Francisco to New York	40 ms	4 years
Internet: San Francisco to United Kingdom	81 ms	8 years
Internet: San Francisco to Australia	183 ms	19 years
TCP packet retransmit	1-3 s	105-317 years







Flash Memory Choices and Tradeoffs



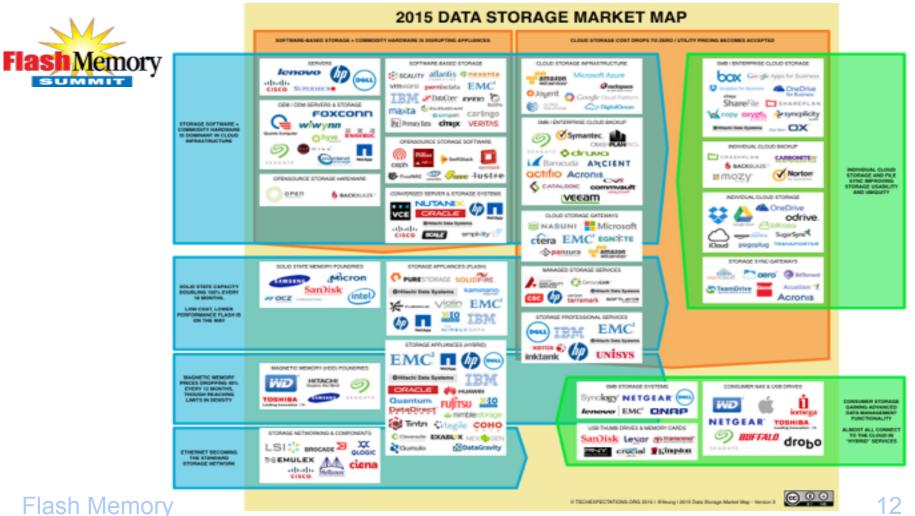
#### Hardware **Architecture**

Software **Development** and Maintenance

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- How big is "Big", and what kind of computer architecture do you need to optimally solve a Big Data problem?
- Big Data applications may have irregular and unpredictable data access patterns
  - Efficient partitioning of data and queries can be challenging
  - Data can change, and structurally vary characteristics over time
- Big Data applications often must solve system problems
  - Coordination of computing resources (Storage, Partitions, Networks, CPU's)
  - People who have deep analytic skills are often required to design system resource strategies, perhaps not the greatest use of their time and expertise



Summit 2015





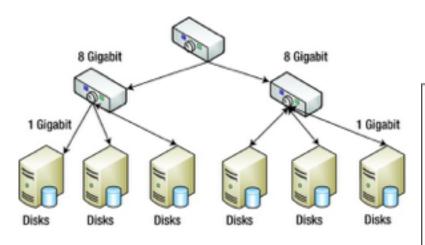


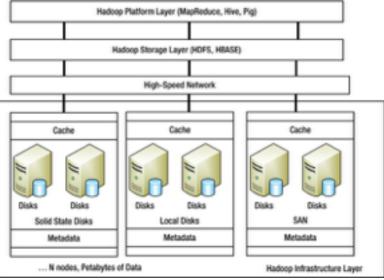
Data Locality is the real problem not, Volume, Velocity or Variety

- The Impact of Data Locality
  - Volume: The further the data has to move the worse your performance will be.
  - Velocity: The more complex the transformations the worse your performance will be.
  - Variety: The more data components you look at the worse your performance will be.



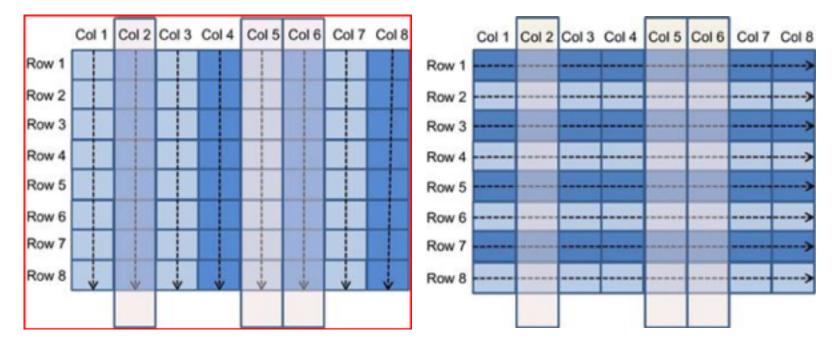
## FlashMemory An Example: Where to optimize?







## Flash Memory A Simple Improvement: Column verses Row data architecture





#### In-Memory Performance

- The world wants data to be in-memory, but hasn't been able to get it.
- Historically the industry has scaled applications to fit on available computer hardware.
- The industry needs to move to scale the hardware to fit the application.

#### Linear System Scalability

- We need systems that enable hardware to scale organically using low cost commodity servers at linear cost, as customer needs evolve.
- We need to enable applications to achieve superior in-memory performance using inexpensive unmodified hardware.

#### Reduced Software Development Costs

- We need to use off-the-shelf, unmodified Linux.
- We require no changes to applications or database software.
- We require systems that optimize automatically, and learn over time.

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- Industry Context
- Addressing the challenge

A proposed solution

Review of the Benefits



- Both scale up and scale out can be expensive
- When all you have is a hammer...
  - Every problem looks like a nail •
- Today we rely almost exclusively on "scale out" systems
  - Because that's the main way we add processors and memory
  - $\rightarrow$  Shard the data, intelligently target the queries time consuming
  - It's not easy to query partitioned databases What is the best way to do it?

    - Moving data is time-consuming
    - And you might have to change it

#### What if you could build systems that "Scale up and Out"?



- Enable application developers to focus time on applications, not the systems required to run them
- Move the coordination and management tasks to the "Operating System" (broadly based)
- Remove the requirement that managing resources be part of the Big Data application

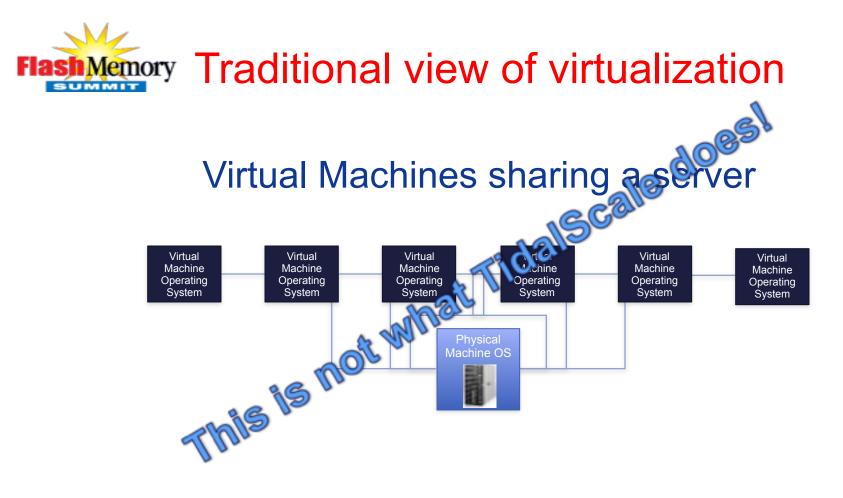




- 45 years ago we figured out how to virtualize memory using the locality principle\*
- Questions:
  - Could locality be applied ubiquitously across our computing infrastructure?
  - How might we apply locality to all compute resource types automatically & dynamically?

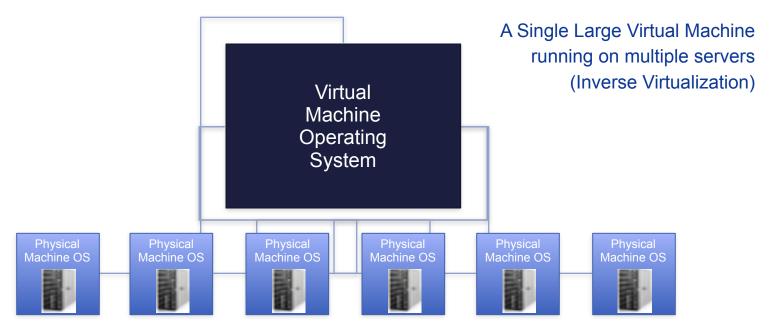


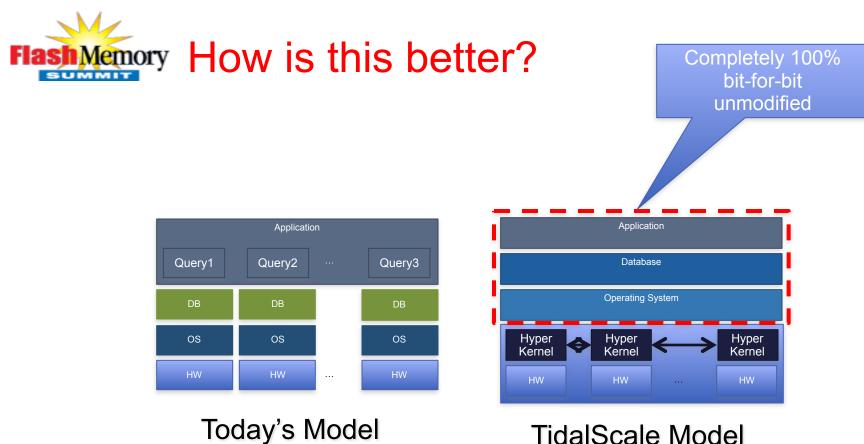
\* P.J. Denning





## TidalScale Provides hardware aggregation





TidalScale Model



### What enables the solution? The Memory Hierarchy in Human Terms

Latency

Scaled

TidalScale Operating Zone	
Memory Cliff	

Event

_	Lyon	Latency	Julieu
70ne	1 CPU Cycle	0.3 ns	1 s
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CIII	Flash SSD (PCI-e) Rotational disk I/O Internet: San Francisco to New York Internet: San Francisco to United Kingdom	4.7 ms 1-10 ms 40 ms 81 ms	5 months

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### Creates a unique scalable solution experience:

#### User experience bare metal

Why?

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#### User experience TidalScale

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#### **Real Screen Shots**

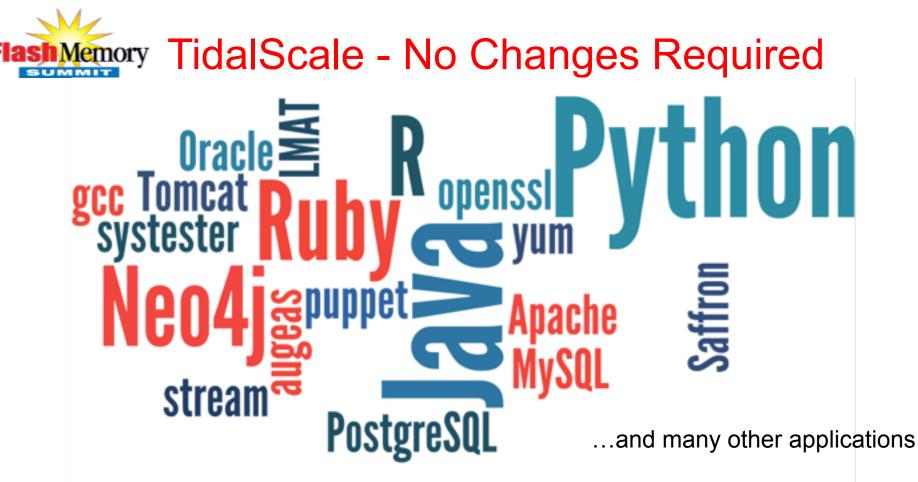


#### System features

Admin node	1
Worker nodes	25
Total Memory	3.2TB
Total Cores:	150
Network	1/10GbE
Storage	FreeNAS, xTB

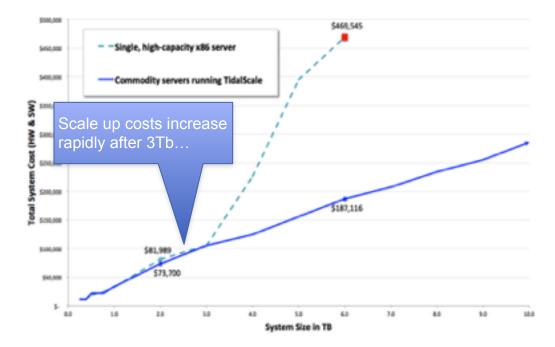
#### Components

1x Admin node (A003) Colfax CX1260i-X6 Haswell, E5-2603V3 6C, 16GB 25x Worker nodes Colfax CX1260i-X6 Haswell, E5-2603V3 6C, 128GB 1x 1G switch 2x 10G switch (S009, S010) Mellanox 1x NAS

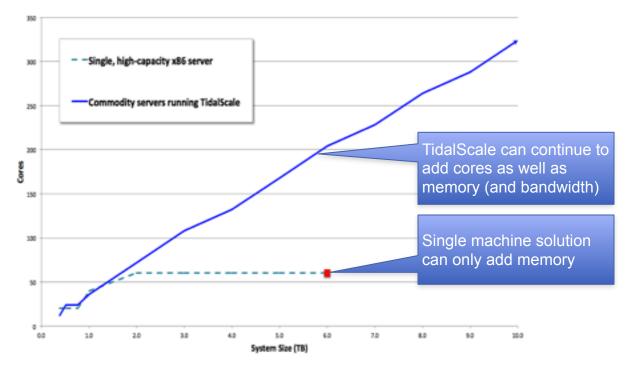


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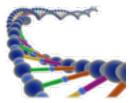


Flash Memory What Problems Benefit?



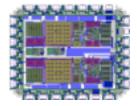
#### Data Mining / Finance

• High Data Volumes, Large Analytics, Risk Analysis, Fraud Detection, Graph Analytics using Alternative Data Sources, Risk modeling, High Frequency trading, Complex Event Modeling





- Next Generation DNA sequencing, Meta-genomic analysis, Finite Element Brain Modeling, Time-Series MRI Neuro-Imaging
- IT / Operational Systems
  - In-House Applications, Web Controllers & Servers, Gateways, Image serving, Ad serving, OLTP, ERP, Business Intelligence





## Flash Memory Scale, Simplify, Optimize, Evolve

#### Scale:

Aggregates compute resources for large scale in-memory analysis and decision support

Scales like a cluster using commodity hardware at linear cost Allow customers to grow gradually as their needs develop

### Simplify:

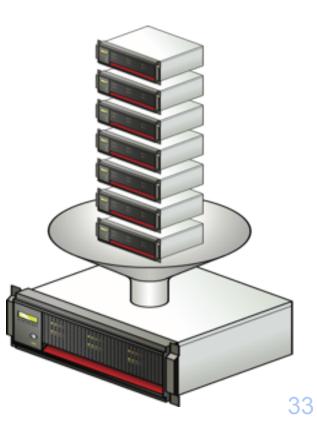
Dramatically simplifies application development No need to distribute work across servers Existing applications run as a single instance, without modification, as if on a highly flexible mainframe

### **Optimize:**

Automatic dynamic hierarchical resource optimization

#### **Evolve:**

Applicable to modern and emerging microprocessors, memories, interconnects, persistent storage & networks





# TidalScale Scale | Simplify | Optimize | Evolve

### Restoring DEVELOPER PRODUCTIVITY THROUGH SIMPLICTY