

Lightweight User-Mode File Systems for Storage Class Memory

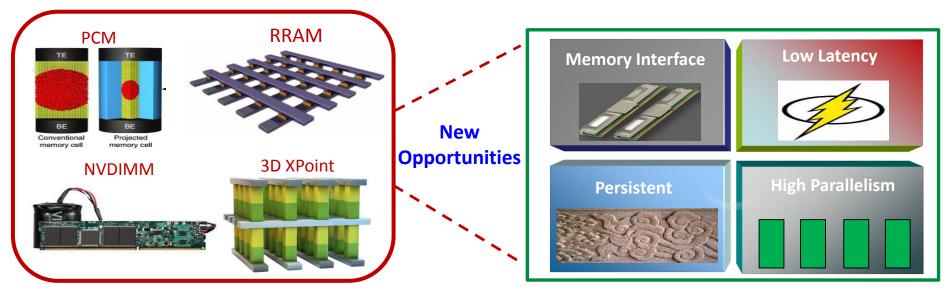
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Flash Memory SCM Era is Coming!



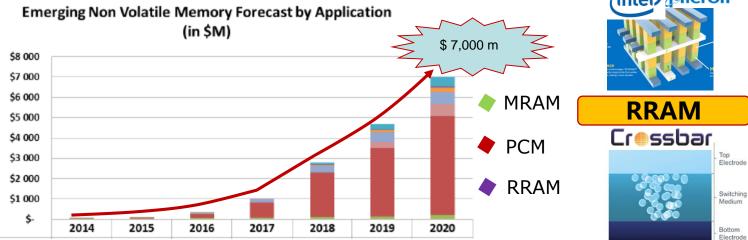
Storage class Memory is the New Storage!





MRAM

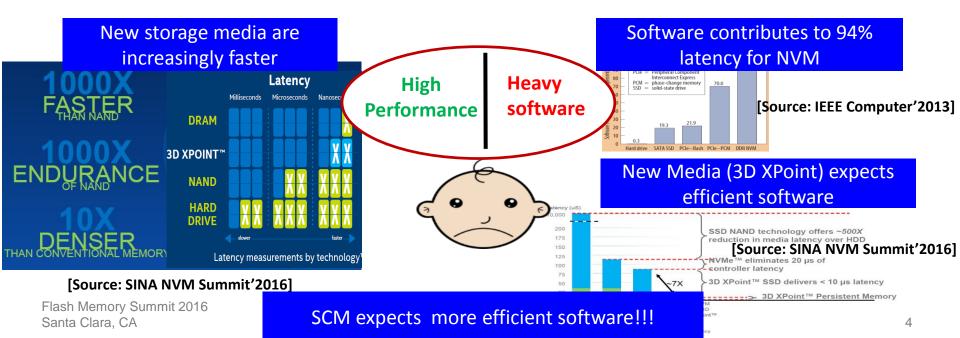
EVERSPIN



[Yole Développement. Emerging Non-volatile Memory. Lyon: Yole Développement; 2015. p. 275.]

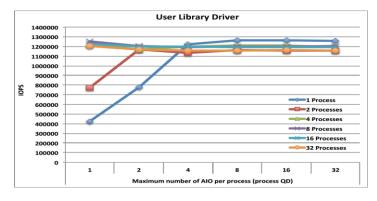


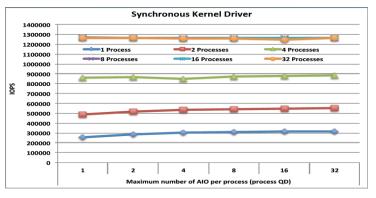
High performance storage like Storage Class Memory (SCM) such as 3D XPoint drives innovation to storage systems





- User Library Driver V.S. Kernel Driver (Performance test for PCM by HGST)
 - Similar peak IOPS, kernel driver demonstrates 8x CPU resource consumption, and
 1.7x access Latency, compared to user driver (3.8 us V.S. 2.2 us)
- Reason: High context switch and system call overhead





Flash Memory Summit 2016 Santa Clara, CA [Source: NVMW'15 Block Device Driver Design for Fast NVM Class Devices]



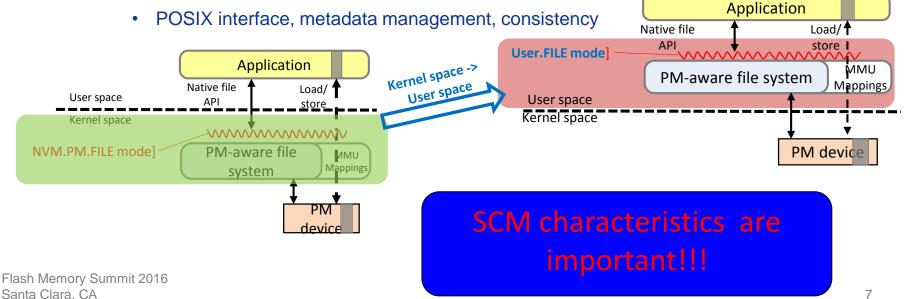
- SCM file systems: BPFS, SCMFS, Aerie, PMFS, NOVA
 - None of exist file system is in user mode meanwhile deep-leverage SCM performance (file indexing latency, parallelism, NUMA-aware)

	User-Mode	Unified File Indexing	Leverage Parallelism	NUMA-Aware Placement
Ramfs	×	×	*	×
BPFS ^[SOSP'09]	×	×	*	×
SCMFS ^[SC'11]	×	×	*	*
Aerie ^[EuroSys'14]	~	×	*	*
PMFS ^[EuroSys'14]	×	×	*	×
EXT4-DAX	×	×	*	×
NOVA ^[FAST'15]	×	×	~	*
LUMFS	\checkmark	~	\sim	\checkmark
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FlashMemory

User-mode File System for SCM

- We develop LUMFS: User-mode system for SCM
 - Bypassing kernel
 - All in user-mode





Outline

- Motivation: User-Mode File Systems
- LUMFS: Lightweight User-Mode File System
- Evaluation
- Conclusion



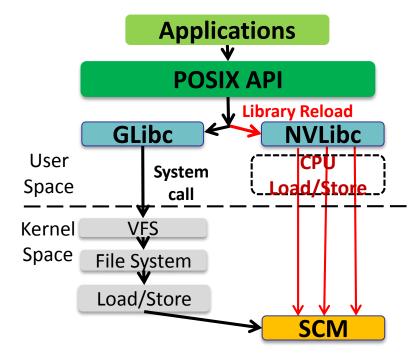
For a user-mode SCM based system that compatible to legacy applications:

- How to handle the POSIX interface?
 - User-mode library "NVLibc"
- How to organize and localize SCM data efficiently?
 - MMU/TLB compatible data management
- How to fully utilize the parallelism of SCM?
 - Resource partitioning
- How to maintain data consistency?
 - Hybrid logging, atomic instructions, CoW



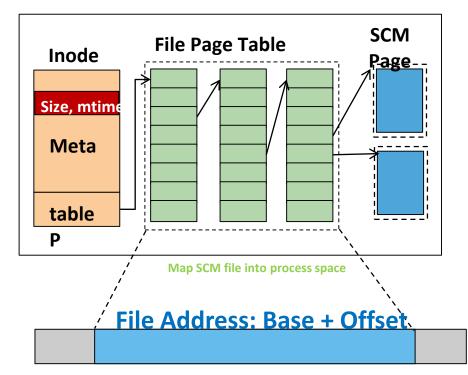
POSIX-compliant Interface

- User-mode Library (NVLibc) provides POSIX interfaces
 - Override GLibc library at runtime by LD_PRELOAD execution environment
 - No modification to source codes
- Use file path to distinguish different FS
 - E.g.,
 - LUMFS/NVlibc: "/mnt/lumfs/test"
 - Kernel/Glibc: otherwise





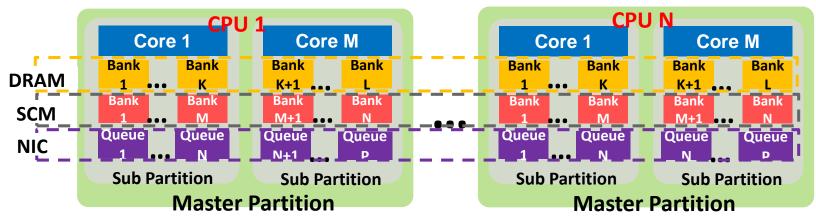
- Key Idea: Unified memory and file indexing
 - All files are mapped into process's virtual address space
 - SCM spaces are organized at page granularity
 - The mapping between file location and SCM pages are managed using page tables, and translated by the MMU



Process Virtual Address Space



Resource Partition



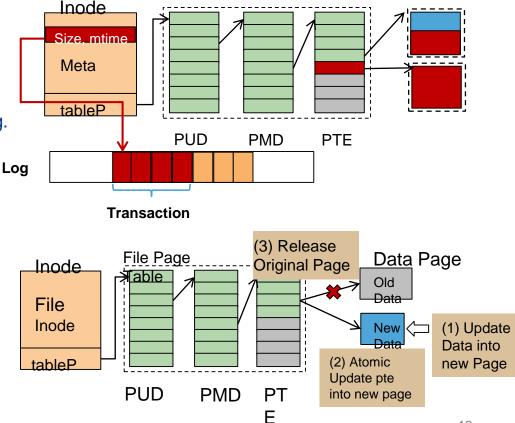
- Resource Partition: to leverage parallelism and reduce contention
 - Master Partition (Per-CPU): Network Card, NUMA Node (SCM & DRAM)
 - Sub Partition (Per-Core): Network card queues, Banks (SCM & DRAM), Caches
- Intelligent Resource Allocation
 - Scatter metadata and file data into multiple sub partitions as possible
 - NUMA-Aware Schedule: Binding thread to local NUMA based on partition of files being accessed
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Metadata & Data Consistency **Flash** Memory

Metadata Consistency

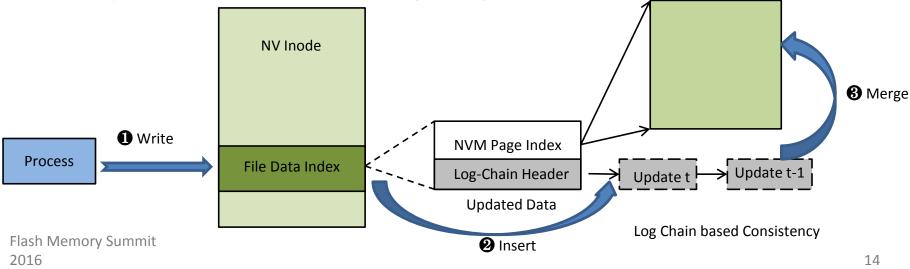
SUMMIT

- Atomic Update for small granularity • update (e.g. size, atime, mtime)
- **Per-partition logging** for large update (e.g. • SCM page allocation)
- Data Consistency
 - Append Write: no worry for data, atomic • update on metadata: (size + mtime)
 - **Copy-On-Write** large granularity data • write, atomic update file page table index
 - **Per-partition data logging: small** granularity data write



Flash Memory Write Consistency : Log Optimization

- Fusion Data and Log optimization
 - Log-Chain to reduce log cost.
 - Updates are linked into Log-Chain (Out of Place update, no log).
 - Read check the part of Log-Chain.
 - Merge Log-Chain into file data ,when large enough



NVM Page Data



- Platform
 - Software: CentOS 6.5, Linux Kernel • 3.5.0, FIO 2.2.10
 - Hardware: •

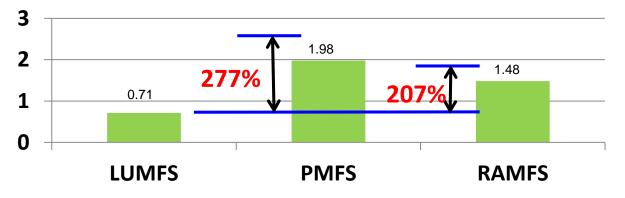
	Server	Huawei RH2288 V3	
Processors	CPUs	2	
	Cores (hyper-thread)	24	
	Model	Intel Xeon E5-2620,	
	IVIOUEI	2.4GHz	
Memory	Channels	8	
	Frequency	DDR4 2133MHz	
	Capacity	384GB	
SCM	Capacity (DRAM Simulated)	192GB	

- **Compared FS**
 - LUMFS : user-mode FS
 - RAMFS: kernel-mode FS, no ٠ consistency
 - PMFS: kernel-mode FS, • consistency

Demo system available at booth 523 (Huawei)

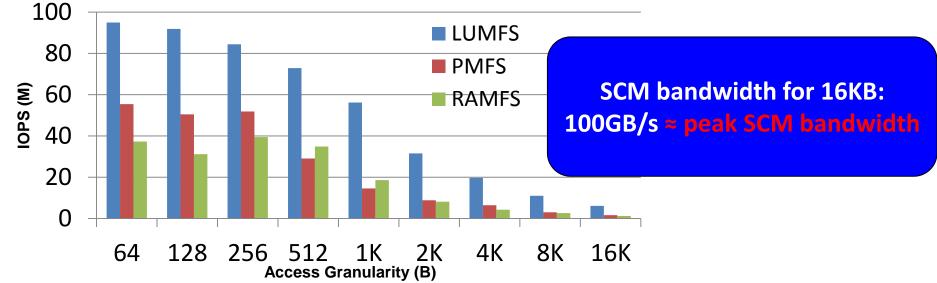


Latency per OP (us)



- LUMFS performs better than kernel-mode FS (PMFS & RAMFS)
 - LUMFS can simplify storage software

Flash Memory FIO performance (random read)



- LUMFS retains its benefits on fully utilizing the parallelism of SCM
 - NUMA, multiple channels, multiple banks

Reduced lock contention through resource partition
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- LUMFS: A User-Mode File System for SCM
 - Manage metadata and data in user mode
 - POSIX-compliant interface support MySQL application
 - **Unified file indexing**: MMU/TLB compliant file page table
 - Thorough resource partitioning leverage SCM parallelism and reduce resource contention
 - Hybrid consistency mechanism for different types of data/metadata
- LUMFS can perform 4KB Random Read/Write IOPS: ~3x PMFS/RAMFS
- It's time for User-Mode File System for SCM



Thank You !

We are hiring! Welcome to join us!

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