# NOVA: A High-Performance, Fault Tolerant File System for Non-Volatile Main Memories

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# **NVMM File System Requirements**

- Legacy File IO Acceleration fast and easy
  - Run existing IO-intensive apps on NVDIMMs
  - "just works"
- Strong atomicity
- DAX Mmap
  - Load-store access & complex data structures
  - More control/responsibility for programs
- Data protection
  - Don't trust the media or other software
  - Support backups
- High performance
  - Otherwise, why bother?





# XFS EXT4 F2FS BTRFS NILFS







# BPFSSCMFSPMFSAerieEXT4-DAXXFS-DAXM1FS







# NOVA







# Files, Atomicity, and Performance







# Log Structured FS For NVMM

- A Nova FS is a tree of logs
  - One log per i-node
  - Logs are not contiguous
- I-nodes point to the head and tail





## Atomicity: Logging for Simple Metadata Operations

- Combines log-structuring, journaling and copy-on-write
- Log-structuring for single log update
  - Write, msync, chmod, etc
  - Lower overhead than journaling and shadow paging





## Atomicity: Lightweight Journaling for Complex Metadata Operations

- Lightweight journaling for update across logs
  - Unlink, rename, etc
  - Journal log tails instead of metadata or data





# **Atomicity: Copy-on-write for file data**

- Copy-on-write for file data
  - Log only contains metadata
  - Log is short





# **Atomicity: DAX**

- Nova does not make atomicity guarantees for DAX mmap()'d data
- msync() works, but it's slow & non-atomic
- The program must ensure consistency
  - ISA Support: CLWB, clflush, etc.
  - Careful programming





# **DRAM Indexes for High Performance**

- Per-inode logging allows for high concurrency
- Split data structure between DRAM and NVMM
  - Persistent log is simple and efficient
  - Volatile tree structure has no consistency overhead





# **Performance and Scalability**

- Put allocator in DRAM
- High scalability
  - Per-CPU NVMM free list, journal and inode table
  - Concurrent transactions and allocation/deallocation





# **Fast garbage collection**

- Log is a linked list
- Log only contains metadata



- Fast GC deletes dead log pages from the linked list
- No copying





# **Thorough garbage collection**

- Starts if valid log entries < 50% log length
- Format a new log and atomically replace the old one





## Recovery

- Normal shutdown recovery:
  - Store allocator state in recovery inode
  - Constant time startup
- Failure recovery:
  - Parallel scan
  - Failure recovery bandwidth:> 400 GB/s







- Intel PM Emulation Platform
  - Emulates different NVM characteristics
  - Emulates clwb/clflush latency
- NOVA provides low latency atomicity



#### **Filebench throughput**

 NOVA achieves high performance with strong data consistency

# DAX, Backup, and Data Protection





# **DAX Challenges**

- DAX mmap() gives the programmer great power
  - Fine-grain updates
  - Pointer-based data structures
  - Custom data layout
- ...along with great responsibilities.
  - Data structure consistency/persistence ordering
  - Memory protection/media error management
- The file system must not interfere.





### **Snapshots for Normal File Access**



# **Memory Ordering With DAX mmap()**

D = 42; V = true;

- D and V live in two pages of a mmap()'d region.
- Recovery invariant: *if V == True, then D is valid*



# **Corrupt Snapshots with DAX-mmap()**

• Recovery invariant: *if V == True, then D is valid* 



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- Incorrect: Naïvely mark pages read-only one-at-a-time

## **Consistent Snapshots with DAX-mmap()**

• Recovery invariant: *if V == True, then D is valid* 

Correct: Block page faults until all pages are read-only



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## **Performance impact of snapshots**

- Normal execution vs. taking snapshots every 10s
  - Negligible performance loss through read()/write()
  - Average performance loss 6.2% through mmap()





# **Protecting Metadata**



# **NVMM Failure Modes: Media Failures**

- Detectable & correctable
  - Transparent to software
- Detectable & uncorrectable
  - Affect a contiguous range of data
  - Raise machine check exception (MCE)
- Undetectable
  - May consume corrupted data
- Software scribbles
  - Kernel bugs or own bugs
  - Transparent to hardware





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  - NVMM file systems are highly vulnerable





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# **NOVA Metadata Protection**

- Replicate everything
  - Inodes
  - Logs
  - Superblock
  - ...
- CRC32 Checksums everywhere





# **Defense Against Scribbles**

- Tolerating Larger Scribbles
  - Allocate replicas far from one another
  - Can tolerate arbitrarily large scribbles to metadata.
- Preventing scribbles
  - Mark all NVMM as read-only
  - Disable CPU write protection while accessing NVMM



# **Protecting Data**





## **NOVA Data Protection**

- Divide 4KB blocks into 512-byte stripes
- Compute a RAID 5-style parity stripe
- Compute and replicate checksums for each stripe



# Data Protection With DAX mmap()

- File systems cannot efficiently protect mmap'ed data, since stores are invisible
- NOVA's data protection contract:

NOVA protects pages from media errors and scribbles iff they are not mmap()'d for writing.





# File data protection with DAX-mmap

• NOVA logs mmap() operations





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# File data protection with DAX-mmap

• On unmap and during recovery, NOVA restores protection





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# Performance





### **Storage Utilization**







#### **Performance Cost of Data Integrity**

# Conclusion

- Existing file systems do not meet the requirements of applications on NVMM file systems
- NOVA's multi-log design achieves high performance and strong consistency
- NOVA's data protection features ensure data integrity
- NOVA outperforms existing file systems while providing stronger consistency and data protection guarantees



# NOVA is open source. We are preparing it for addition to Linux.

To help or try it out: https://github.com/ NVSL/linux-nova









# Thanks!

